

# Spectral limit theorems for submatrices and products of projections

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# **I. Introduction**

# Eigenvalues Under Compression

$\mathbf{A}, \mathbf{B} \in \mathbb{C}^{N \times N}$  orthogonal projections of **linear rank**:

$$\frac{1}{N} \text{Tr}(\mathbf{A}) \approx \alpha \in (0, 1),$$
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This specifies the **empirical spectral distributions**:

$$\text{e.s.d. of } \mathbf{A} \approx (1 - \alpha)\delta_0 + \alpha\delta_1 = \text{Ber}(\alpha),$$

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How does **compressing**  $\mathbf{B}$  by  $\mathbf{A}$  change the eigenvalues?

$$\text{e.s.d. of } \mathbf{ABA} \approx ?$$

# Geometric Interpretation

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$$\cos(\theta_k) := \max_{\substack{\mathbf{a}_k \in U \\ \|\mathbf{a}_k\|=1 \\ \langle \mathbf{a}_i, \mathbf{a}_k \rangle = 0 \text{ for } 1 \leq i < k}} \max_{\substack{\mathbf{b}_k \in V \\ \|\mathbf{b}_k\|=1 \\ \langle \mathbf{b}_j, \mathbf{b}_k \rangle = 0 \text{ for } 1 \leq j < k}} \langle \mathbf{a}_k, \mathbf{b}_k \rangle.$$

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Populate columns of  $\mathbf{U}, \mathbf{V}$  with orthonormal bases of  $U, V$ , so  $\mathbf{A} = \mathbf{U}\mathbf{U}^*$ ,  $\mathbf{B} = \mathbf{V}\mathbf{V}^*$ . Then:

$$\begin{aligned} \cos(\theta_k) &= k\text{th singular value of } \mathbf{U}^* \mathbf{V} \\ &= (k\text{th eigenvalue of } \mathbf{U}^* \mathbf{V} \mathbf{V}^* \mathbf{U})^{1/2} \\ &= (k\text{th eigenvalue of } \mathbf{U} \mathbf{U}^* \mathbf{V} \mathbf{V}^* \mathbf{U} \mathbf{U}^*)^{1/2} \\ &= (k\text{th eigenvalue of } \underbrace{\mathbf{A} \mathbf{B} \mathbf{A}}_{\text{"angle operator"}})^{1/2}. \end{aligned}$$

# Application: Submatrices

Special case:  $A$  is diagonal, a **coordinate projection**:

$$A = \begin{bmatrix} 0 & & & & \\ & 1 & & & \\ & & 1 & & \\ & & & 0 & \\ & & & & 1 \end{bmatrix}$$

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Then,  $ABA$  extracts a **submatrix**:

$$A \begin{bmatrix} \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \end{bmatrix} A = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & \cdot & \cdot & 0 & \cdot \\ 0 & \cdot & \cdot & 0 & \cdot \\ 0 & 0 & 0 & 0 & 0 \\ 0 & \cdot & \cdot & 0 & \cdot \end{bmatrix}$$

## Submatrices $\rightsquigarrow$ Induced Subgraphs

A **strongly  $d$ -regular graph** has only three eigenvalues: the “trivial”  $d$  eigenvalue, and two with large multiplicity:

$$\begin{aligned}\mathbf{G} &= \frac{d}{N} \mathbf{1}\mathbf{1}^* + \lambda_1 \mathbf{B}_1 + \lambda_2 \mathbf{B}_2 \\ &= \frac{d}{N} \mathbf{1}\mathbf{1}^* + \lambda_1 \mathbf{B}_1 + \lambda_2 \left( \mathbf{I} - \mathbf{B}_1 - \frac{1}{N} \mathbf{1}\mathbf{1}^* \right) \\ &= (\lambda_1 - \lambda_2) \mathbf{B}_1 + \text{simple adjustment.}\end{aligned}$$

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If  $(\mathbf{A}_S)_{ii} = \mathbb{1}\{i \in S\}$  for  $S \subseteq [N]$ , then

$\mathbf{A}_S \mathbf{G} \mathbf{A}_S$  = adjacency matrix of induced subgraph on  $S$ ,

and we can understand the spectrum via  $\mathbf{A}_S \mathbf{B}_1 \mathbf{A}_S$ .

# Submatrices $\rightsquigarrow$ Restricted Isometry Property

In **compressed sensing**, want  $\mathbf{v}_1, \dots, \mathbf{v}_N \in \mathbb{C}^M$  with  $N \gg M$  so that any small subset is close to orthonormal.

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$$\mathbf{V} = \begin{bmatrix} | & & | \\ \mathbf{v}_1 & \cdots & \mathbf{v}_N \\ | & & | \end{bmatrix},$$

$$\begin{aligned} k\text{-RIP constant} &= \max_{i_1, \dots, i_k \in [N]} \left\| \left( \langle \mathbf{v}_{i_a}, \mathbf{v}_{i_b} \rangle \right)_{a,b=1}^k - \mathbf{I}_k \right\| \\ &= \max_{S \subset [N], |S|=k} \left\| \mathbf{A}_S \mathbf{V}^* \mathbf{V} \mathbf{A}_S - \mathbf{I}_k \oplus \mathbf{0} \right\|. \end{aligned}$$

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The  $\mathbf{v}_i$  are a **tight frame** if  $\sum_{i=1}^N \mathbf{v}_i \mathbf{v}_i^* = \mathbf{V} \mathbf{V}^* = c \mathbf{I}_M$ .

If so,  $\mathbf{V}^* \mathbf{V} = c \mathbf{B}$  is a rescaled projection, so this is a question about the eigenvalues (over all  $S$ ) of  $\mathbf{A}_S \mathbf{B} \mathbf{A}_S$ .

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**Summary:** If  $U$  and  $V$  are in “sufficiently general position,” then the eigenvalues of  $ABA$  follow the universal **Wachter MANOVA distribution** with density:

$$\begin{aligned}d\mu_{\alpha,\beta}(x) &= \frac{\sqrt{(r_+ - x)(x - r_-)}}{2\pi x(1 - x)} \mathbb{1}_{[r_-, r_+]}(x) dx \\ &\quad + \max\{1 - \beta, 1 - \alpha\} \delta_0(x) \\ &\quad + \max\{\beta - (1 - \alpha), 0\} \delta_1(x), \\ r_{\pm} &= \alpha + \beta - 2\alpha\beta \pm 2\sqrt{\alpha(1 - \alpha)\beta(1 - \beta)} \\ &= \left(\sqrt{\alpha(1 - \beta)} \pm \sqrt{\beta(1 - \alpha)}\right)^2 \in (0, 1).\end{aligned}$$

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**Interpretation:** Eigenvalues  $(1 - \beta)\delta_0 + \beta\delta_1$  of  $B$  are **smoothed**, **0 atom increases**, **1 atom decreases**.

## **II. Empirical Spectral Distribution**

# Free Probability Perspective [Voiculescu '90s]

Consider **sequences of random orthogonal projections**  
 $\mathbf{A}^{(N)}, \mathbf{B}^{(N)} \in \mathbb{C}^N$  with

$$\lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \operatorname{Tr} \mathbf{A}^{(N)} = \alpha, \quad \lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \operatorname{Tr} \mathbf{B}^{(N)} = \beta.$$

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Equivalent to convergence of e.s.d.'s to  $\operatorname{Ber}(\alpha), \operatorname{Ber}(\beta)$ .

Free probability  $\rightsquigarrow$  if  $(\mathbf{A}^{(N)}, \mathbf{B}^{(N)})$  **asymptotically free**, have convergence (in moments) of e.s.d. of  $\mathbf{A}^{(N)} \mathbf{B}^{(N)} \mathbf{A}^{(N)}$  to

$$\operatorname{Ber}(\alpha) \boxtimes \operatorname{Ber}(\beta) = \mu_{\alpha, \beta}.$$

To establish weak convergence, suffices to establish asymptotic freeness.

# Asymptotic Freeness for Projections

Usual definition: for all  $s_1, t_1, \dots, s_k, t_k \geq 1$ , let

$$a_i := \lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \operatorname{Tr} \mathbf{A}^{(N)s_i}, \quad b_i := \lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \operatorname{Tr} \mathbf{B}^{(N)t_i}.$$

Then, asymptotic freeness  $\Leftrightarrow$  for any such choice,

$$\lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \operatorname{Tr} \prod_{i=1}^k (\mathbf{A}^{(N)s_i} - a_i \mathbf{I}_N) (\mathbf{B}^{(N)t_i} - b_i \mathbf{I}_N) = 0.$$

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But for projections, by idempotence, enough to analyze **one-parameter family of traces**  $s_1 = t_1 = \dots = s_k = t_k = 1$ :

$$\lim_{N \rightarrow \infty} \frac{1}{N} \mathbb{E} \operatorname{Tr} ((\mathbf{A}^{(N)} - \alpha \mathbf{I}_N) (\mathbf{B}^{(N)} - \beta \mathbf{I}_N))^k = 0.$$

# Main Theorem 1 [K '23]

$\mathbf{A}^{(N)}, \mathbf{B}^{(N)} \in \mathbb{C}^{N \times N}$  random orthogonal projections with

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Then, we have **convergence in moments**: for all  $k \geq 1$ ,

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Remarks:

- Straightforward application of free probability tools.
- Extended to weak convergence in probability or a.s.

# Application: Random Subsets of Frames

Main Theorem 1 applies with:

- $\mathbf{A}^{(N)}$  diagonal,  $A_{ii} \stackrel{\text{iid}}{\sim} \text{Ber}(\alpha)$ ,
- $\mathbf{B}^{(N)} = \frac{M}{N} \mathbf{V}^{(N)*} \mathbf{V}^{(N)}$  for  $\mathbf{V}^{(N)} = [\mathbf{v}_1 \cdots \mathbf{v}_N] \in \mathbb{C}^{M \times N}$   
(deterministic!) tight frames having  $\frac{M}{N} \rightarrow \beta \in (0, 1)$  and

$$\max_{i,j \in [N]} |\langle \mathbf{v}_i, \mathbf{v}_j \rangle - \mathbb{1}\{i = j\}| \leq N^{-1/2+o(1)}.$$

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Answers signal processing and combinatorics questions:

- Proves conjecture of [Haikin, Zamir, Gavish '17]
- Simplifies [Mixon, Magsino, Parshall '21] (Paley frames)
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...with easy proof!

# Proof Sketch

Just expand and bound naively:

$$\begin{aligned} & \frac{1}{N} |\mathbb{E} \operatorname{Tr} ((\mathbf{A} - \alpha \mathbf{I})(\mathbf{B} - \beta \mathbf{I}))^k| \\ & \leq \frac{1}{N} \sum_{i_1, \dots, i_k=1}^N |\mathbb{E}[(A_{i_1 i_1} - \alpha) \cdots (A_{i_k i_k} - \alpha)]| \\ & \quad |B_{i_1 i_2} - \beta \mathbb{1}\{i_1 = i_2\}| \cdots |B_{i_k i_1} - \beta \mathbb{1}\{i_k = i_1\}| \\ & \lesssim_k N^{-1} \cdot \underbrace{N^{k/2}}_{\# \text{ non-zero terms}} \cdot N^{-k/2+o(1)} \\ & \rightarrow 0. \end{aligned}$$

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Previous work finds a “main term” and “error term” in  $\frac{1}{N} \mathbb{E} \operatorname{Tr}(\mathbf{A}\mathbf{B}\mathbf{A})^k$  directly, redoing free probability by hand.

## Aside: Free Probability with “Less Randomness”

Classical results (Voiculescu et al.) treat **“very random”** models:  $A, B$  **unitarily invariant** projections, i.e., to uniformly random subspaces of  $\mathbb{C}^N$ .

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Companion work [K '23]: some cases where the same limits hold for **completely deterministic** models, e.g., built on the Paley frames and Paley graph of number theory.

# Asymptotic Freeness in Paley Graphs

$G_p$  a graph on vertices  $\mathbb{Z}/p\mathbb{Z}$  (with  $p \equiv 1 \pmod{4}$ ) with  $i \sim j$  iff  $j - i$  is a **square** mod  $p$  (for some  $x \neq 0$ ,  $j - i \equiv x^2$ ).

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**Theorem:** [K '23] The associated projection  $B$  is asymptotically free of coordinate projections  $A_S$  for any

$$S = \{j : j \sim i \text{ in } G_p\} = \text{“neighborhood of } i\text{”}.$$

So, e.s.d.  $\rightarrow \mu_{1/2,1/2} =$  **arcsine law**. (Partial generalization to other “low-degree” sets  $S$ .)

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**Moral:** Freeness does not require probability!

## **III. The Largest Eigenvalue**

# Conjecture and Difficulties

Natural extension of previous results:

**Edge Conjecture:** For many  $A^{(N)}, B^{(N)}$  with e.s.d. of  $A^{(N)} B^{(N)} A^{(N)}$  converging weakly to  $\mu_{\alpha, \beta}$ ,

$$\lambda_{\max}(A^{(N)} B^{(N)} A^{(N)}) \xrightarrow{(p)} \text{edge}(\alpha, \beta) := \text{right edge of } \mu_{\alpha, \beta}.$$

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As in models from classical random matrix theory (e.g., Wigner and Wishart), the difficulty is in controlling moments of **diverging order**:

$$\mathbb{E} \text{Tr}(\mathbf{A}^{(N)} \mathbf{B}^{(N)} \mathbf{A}^{(N)})^k \stackrel{?}{\leq} (\text{edge}(\alpha, \beta) + o(1))^k$$

for  $k \gg \log(N)$ .

## Main Theorem 2

For  $\mathbf{A}^{(N)}, \mathbf{B}^{(N)}$  as in Main Theorem 1, suppose additionally that **one of  $\alpha$  or  $\beta$  is  $\frac{1}{2}$**  and that for  $k = k(N) \gg \log(N)$ ,

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**Conjecture:** **Extra condition on  $\alpha, \beta$**  not necessary.

Seems just a technical challenge—stay tuned for details.

## Main Theorem 2: Intuition

Why do these normalizations appear?

$$\hat{\mathbf{A}} := (\mathbf{A}^{(N)} - \alpha \mathbf{I}_N) / \sqrt{\alpha(1 - \alpha)}, \quad \hat{\mathbf{B}} := (\mathbf{B}^{(N)} - \beta \mathbf{I}_N) / \sqrt{\beta(1 - \beta)}$$

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**Normalization whitens the spectrum:** if  $\lambda_i(\mathbf{A}^{(N)}) \stackrel{\text{iid}}{\sim} \text{Ber}(\alpha)$  and  $\lambda_i(\mathbf{B}^{(N)}) \stackrel{\text{iid}}{\sim} \text{Ber}(\beta)$ , then

$$\begin{aligned}\mathbb{E}\lambda_i(\hat{\mathbf{A}}) &= \mathbb{E}\lambda_i(\hat{\mathbf{B}}) = 0, \\ \mathbb{E}\lambda_i(\hat{\mathbf{A}})^2 &= \mathbb{E}\lambda_i(\hat{\mathbf{B}})^2 = 1.\end{aligned}$$

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$\rightsquigarrow \hat{\mathbf{A}}, \hat{\mathbf{B}}$  are **orthogonal in expectation**. Actually orthogonal matrices would indeed satisfy

$$|\text{Tr}(\hat{\mathbf{A}}\hat{\mathbf{B}})^k| \leq N = \exp(o(k)).$$

## Proof Ideas for Main Theorem 2

Exact relation between centered and uncentered moments by “trace rewriting” using idempotence:

$$\begin{aligned} & \text{Tr}((\mathbf{A} - \alpha \mathbf{I}_N)(\mathbf{B} - \beta \mathbf{I}_N))^k \\ &= w \text{Tr}(\mathbf{I}_N) + x \text{Tr}(\mathbf{A}) + y \text{Tr}(\mathbf{B}) + \sum_{\ell=1}^k z_\ell \text{Tr}(\mathbf{ABA})^\ell. \end{aligned}$$

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Inverting this, the “main term” is the limiting MANOVA moment:

$$\begin{aligned} \text{Tr}(\mathbf{ABA})^k &= N \mathbb{E}_{\lambda \sim \mu_{\alpha, \beta}} \lambda^k + x' (\text{Tr}(\mathbf{A}) - \alpha N) + y' (\text{Tr}(\mathbf{B}) - \beta N) \\ &\quad + \sum_{\ell=1}^k z'_{\ell} \text{Tr}((\mathbf{A} - \alpha \mathbf{I}_N)(\mathbf{B} - \beta \mathbf{I}_N))^k. \end{aligned}$$

# Proof Ideas for Main Theorem 2

First result: an **explicit recursion** for the moment errors

$$\begin{aligned}\Delta_k &= \text{Tr}(\mathbf{A}\mathbf{B}\mathbf{A})^k - N \mathbb{E}_{\lambda \sim \mu_{\alpha, \beta}} \lambda^k \\ &\approx \sum_{\ell=1}^{k-1} c_{k, \ell} \Delta_\ell + c'_k \text{Tr}((\mathbf{A} - \alpha \mathbf{I}_N)(\mathbf{B} - \beta \mathbf{I}_N))^k.\end{aligned}$$

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Still not easy to analyze for large  $k$ ...

**Saving grace:** if  $\alpha$  or  $\beta$  is  $\frac{1}{2}$ , can solve this recursion in **closed form!**

Identify **Riordan arrays** in recursion: triangular matrices with special generating functions allowing formal inversion.

## “Application:” New Proof for Invariant Model

New, arguably **more robust**, proof of edge limit theorem for  $\mathbf{A}^{(N)}, \mathbf{B}^{(N)}$  unitarily invariant projections.

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$\mathbf{B} = \mathbf{UDU}^*$ .

Proof: expand; use non-asymptotic bounds for Weingarten function [Collins, Matsumoto '17] to control moments

$$\mathbb{E}[U_{i_1 j_1} \cdots U_{i_k j_k} \overline{U_{i'_1 j'_1}} \cdots \overline{U_{i'_k j'_k}}].$$

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No expectation over  $B$ ; expanding traces leads to trying to analyze “**graphical moments**” of  $B$ :

From a graph  $G = ([k], E)$  and  $B \in \mathbb{R}_{\text{sym}}^{N \times N}$ , compute

$$\sum_{i_1, \dots, i_k \in [N] \text{ distinct}} \prod_{\{a, b\} \in E} B_{i_a i_b}.$$

Scalar version of **graph matrices** appearing in literature on **sum-of-squares optimization**.

Techniques to analyze for **general  $G$  and deterministic  $B$** ?

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**Theorem:** [K '23] If Edge Conjecture holds for Paley graph construction, then largest clique in  $G_p$  is  $\leq o(\sqrt{p})$ .

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- Other universal features:
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  - Rate of convergence of e.s.d.?
- Adapt to analyze RIP (**all** small submatrices)?

**Thank you!**